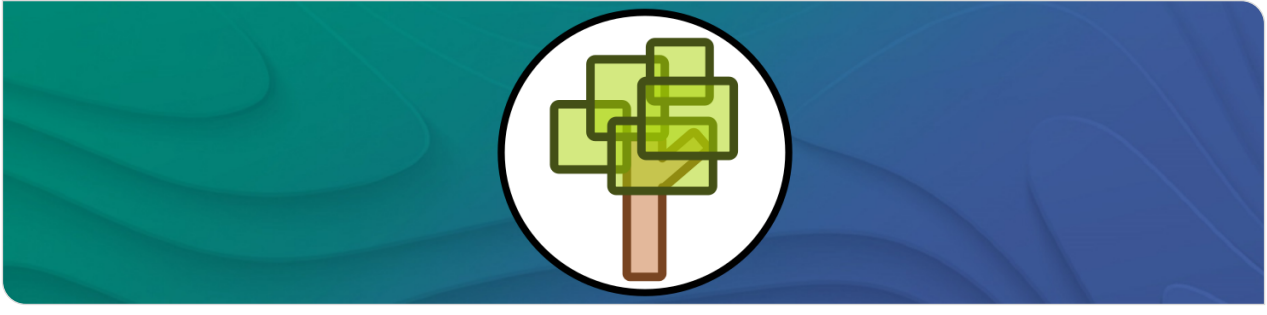
Karlsruhe Institute of Technology

# Faster Block Tree Construction

**31st European Symposium on Algorithms (ESA 2023)**

Dominik Köppl, *Florian Kurpicz*, and Daniel Meyer





# Motivation

- Git repositories
- DNA
- proteins
- user generated content
- XML data
- book collections
- . . .

highly repetitive
&
huge input

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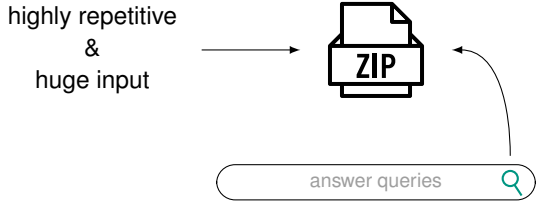

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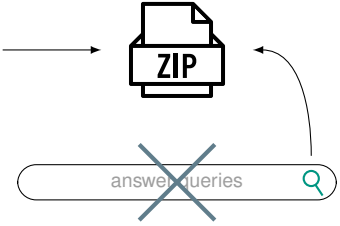

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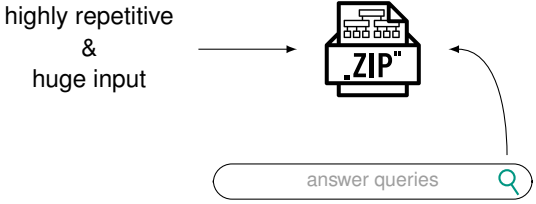

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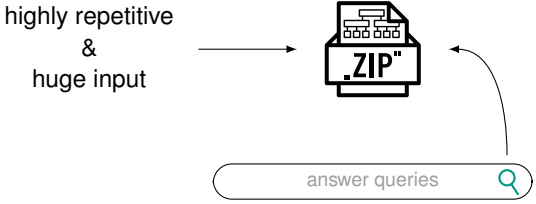


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- rank
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- pattern matching

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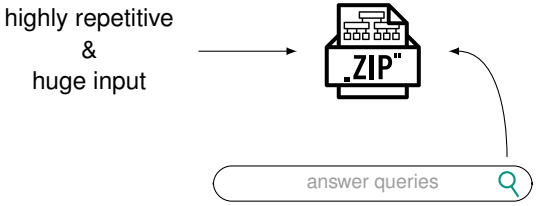


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**Wavelet Tree** (de-facto standard in practice)

- $T \in [1, \sigma]^n$
- access, rank, select: $O(\log \sigma)$ time
- $nH_0(T) + o(n)$ bits space

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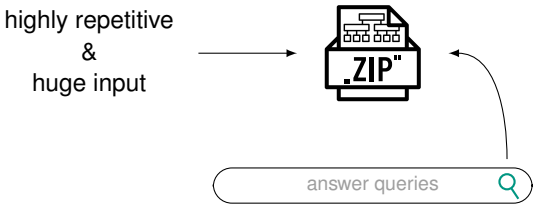


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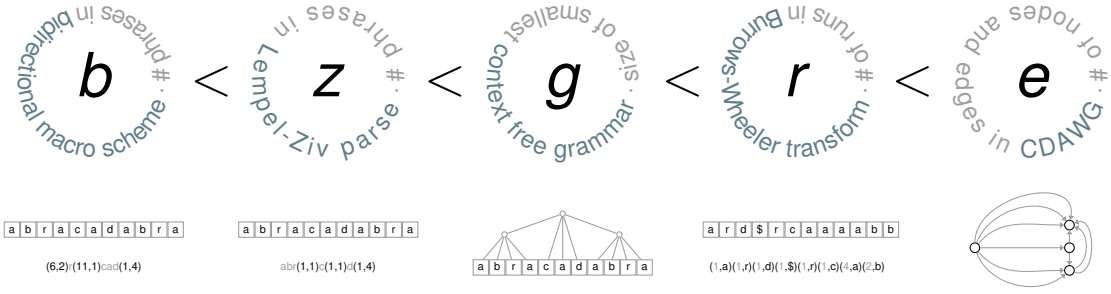
- **blind for repetitions**

# Measures of Repetitiveness (Excerpt)




NP-hard | $O(n)$ time | NP-hard* | $O(n)$ time | $O(n)$ time

* there are good heuristics

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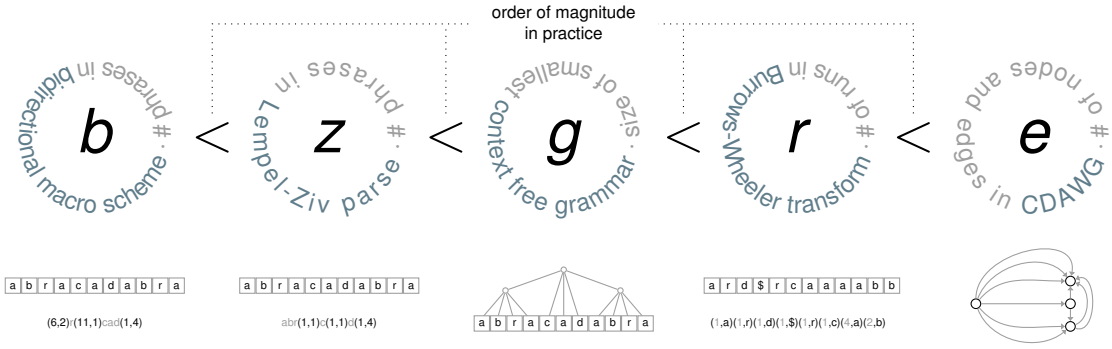


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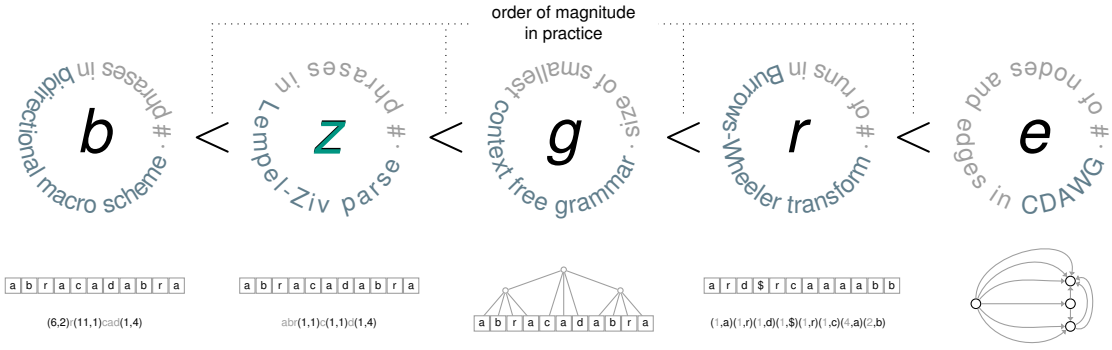


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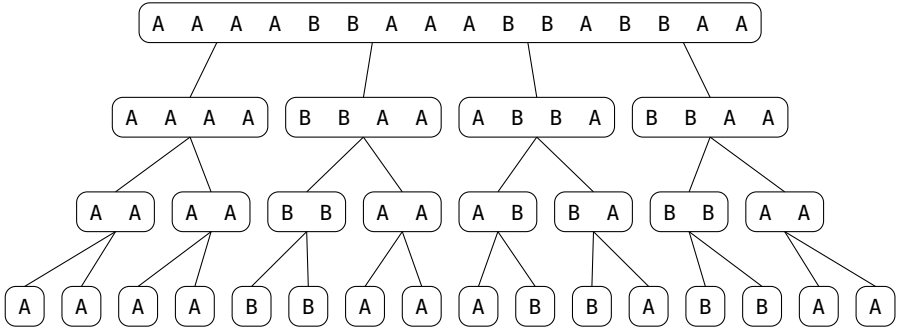
# The Block Tree




- root has degree $s = 4$
- other nodes have degree $\tau = 2$
- all levels (except the first) have $\leq 3z\tau$ blocks

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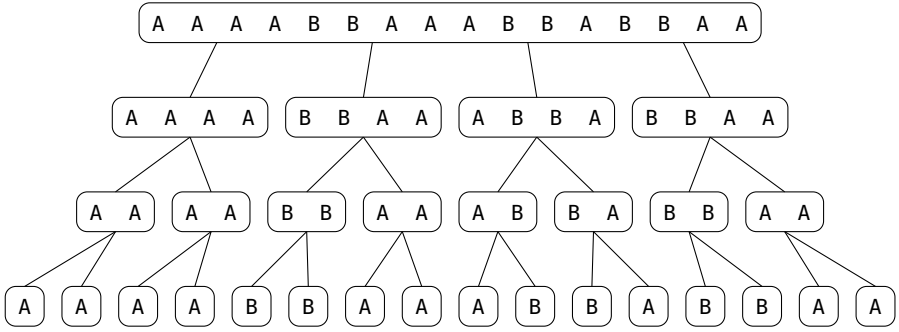


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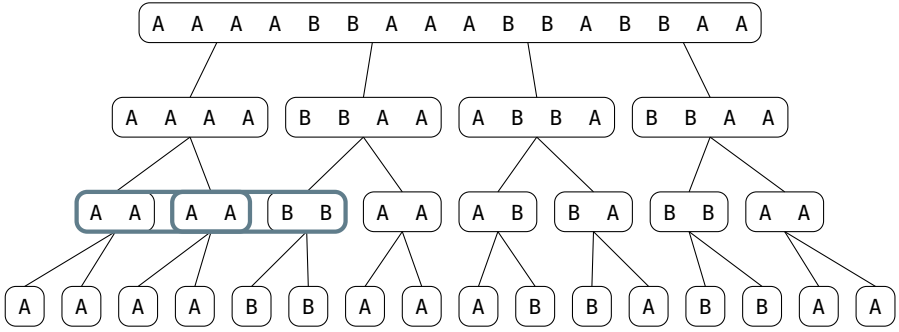


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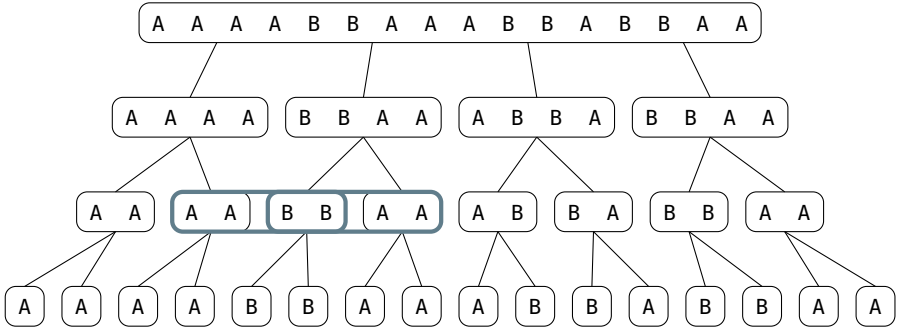


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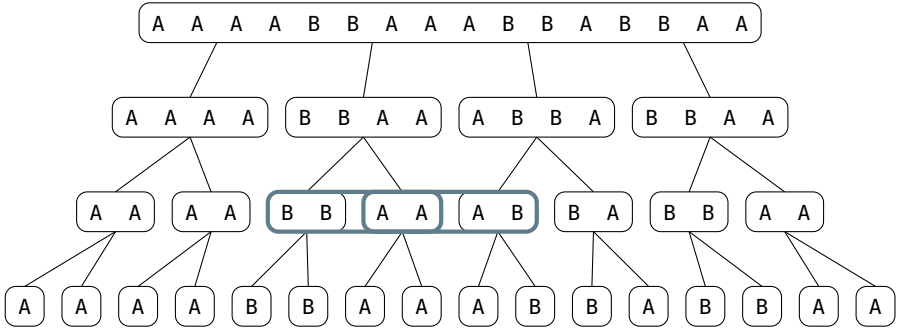


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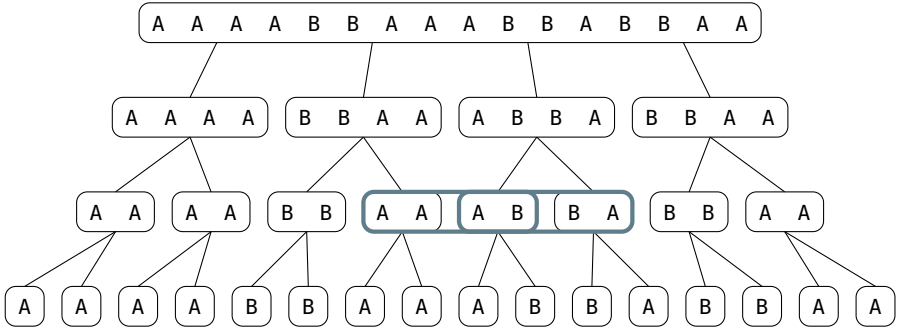


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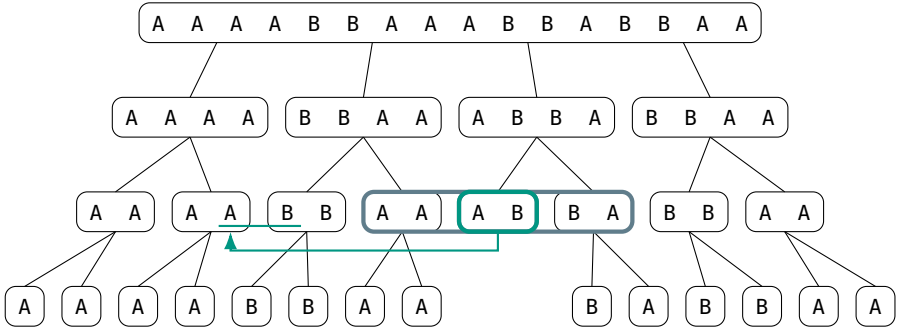


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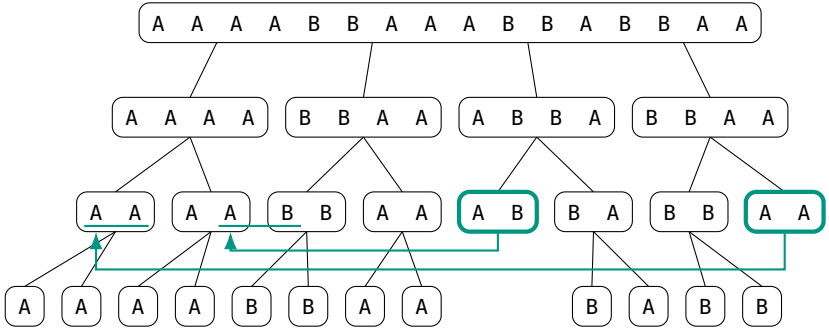


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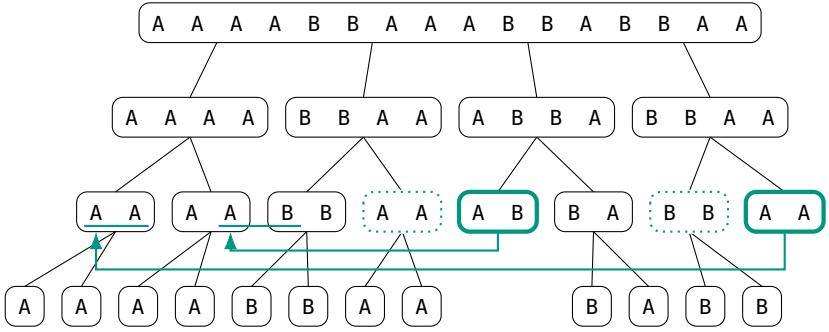


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- keep $B_i$ if consecutive $B_{i-1} \cdot B_i$ or $B_i \cdot B_{i+1}$ are leftmost occurrences
- remove more blocks with pruning

# State-of-Block-Tree-Construction

| Method | Reference | Working Space | Time | Implementation |
|---|---|---|---|---|
| Aho-Corasic | [Bel+21] | $O(n)$ | $O(n(1+\log_\tau(z\tau/s)))$ | no |
| Fingerprints | [Bel+21] | $O(s + z\tau \log_\tau(\frac{n \log \sigma}{s \log n}))$ | $O(n(1+\log_\tau(z\tau/s)))$ expected | yes (slow) |
| LPF Array | [here] | $O(n)$ | $O(n(1+\log_\tau(z\tau/s)))$ | yes (fast) |

# Lempel-Ziv Parse

A A A A B B A A A B B A B B A A

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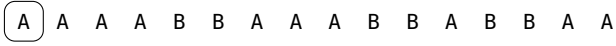

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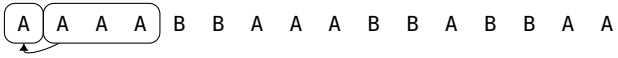

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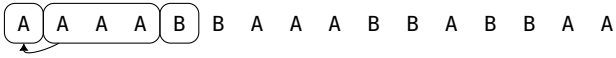

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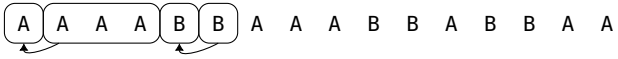

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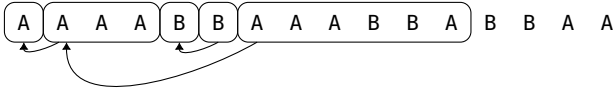

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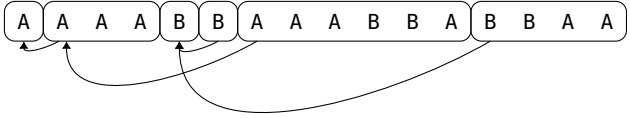

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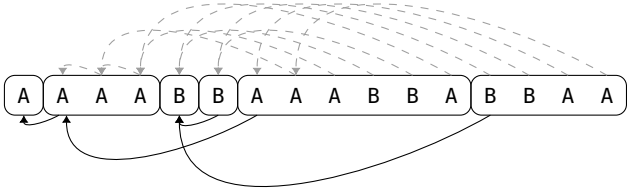

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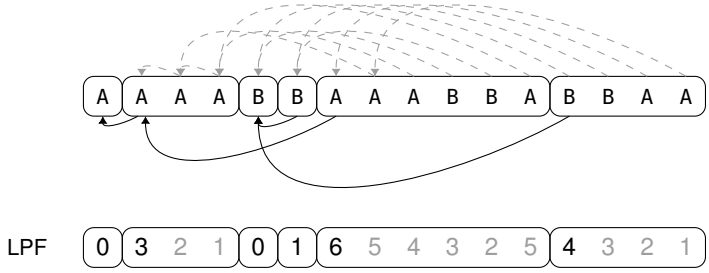

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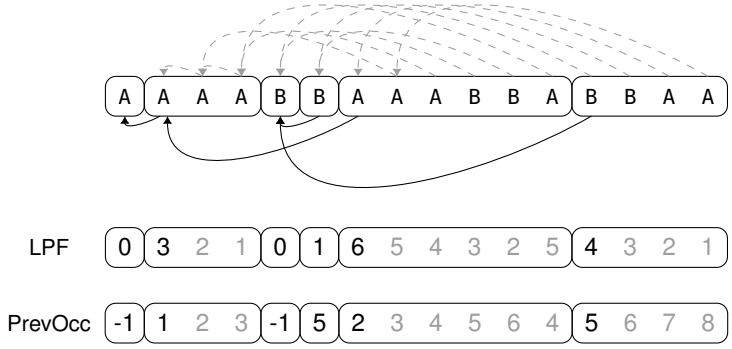

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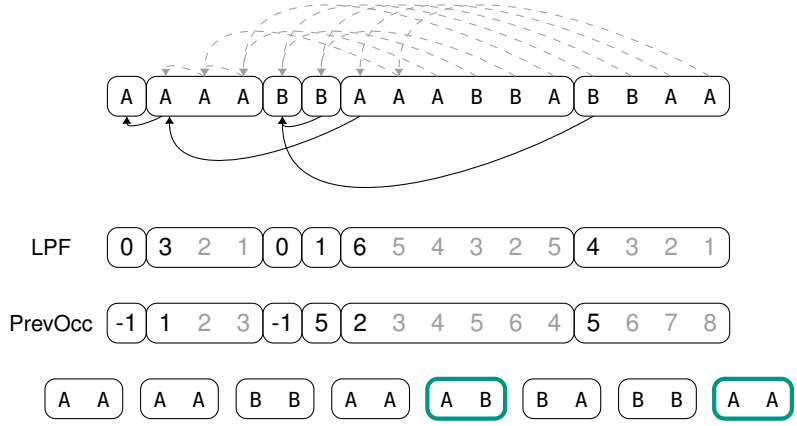

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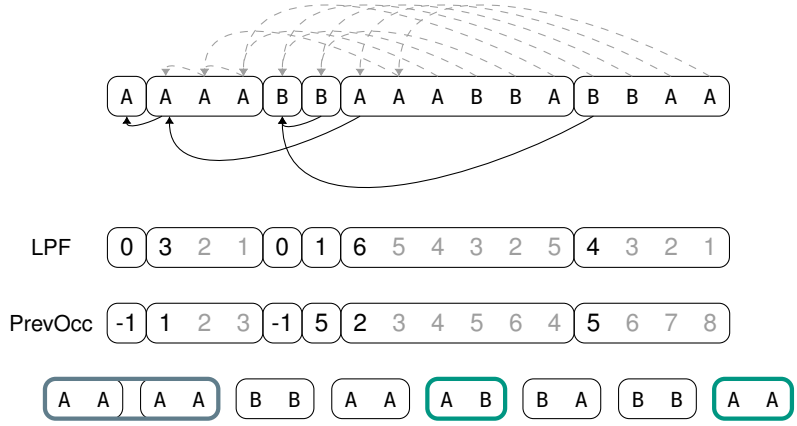

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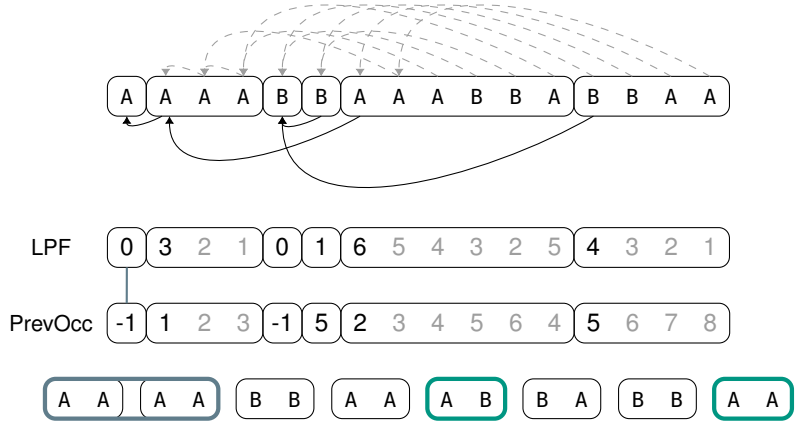

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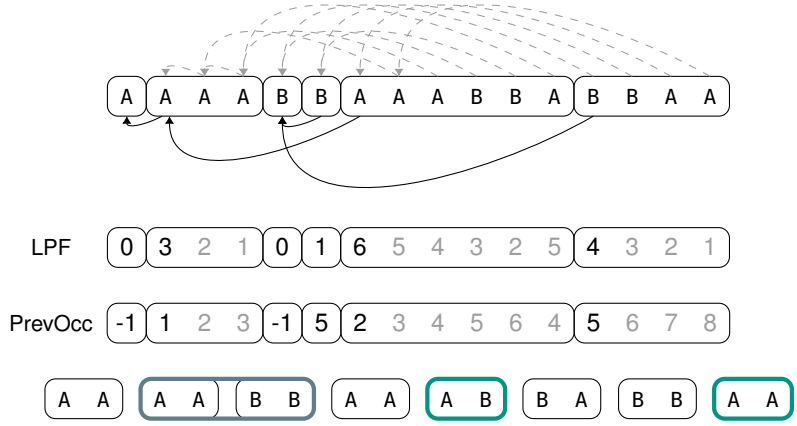

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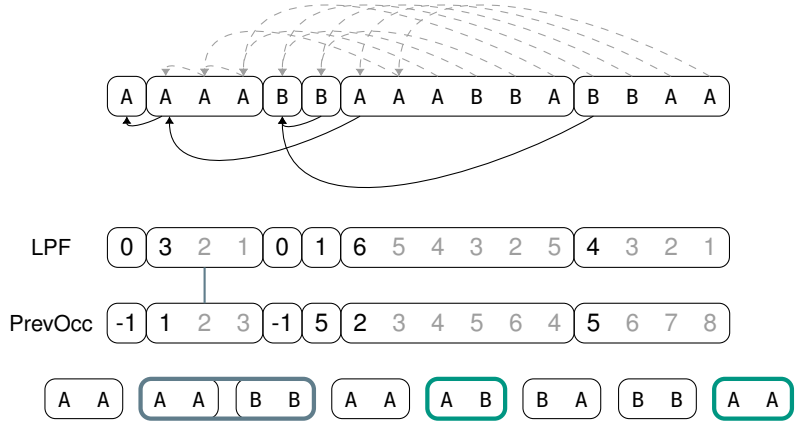

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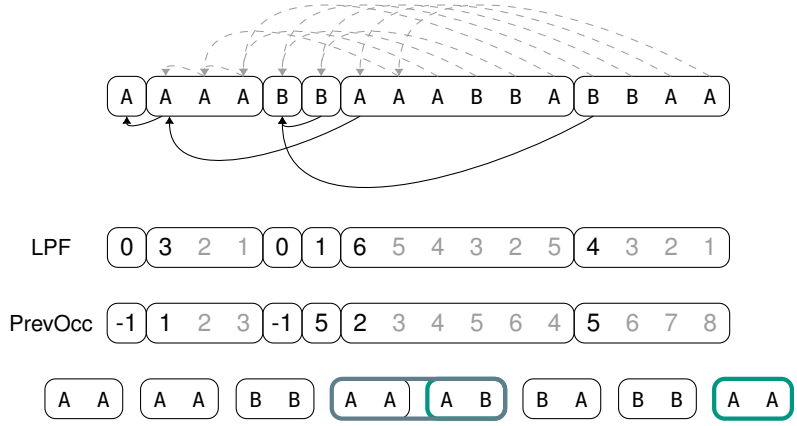

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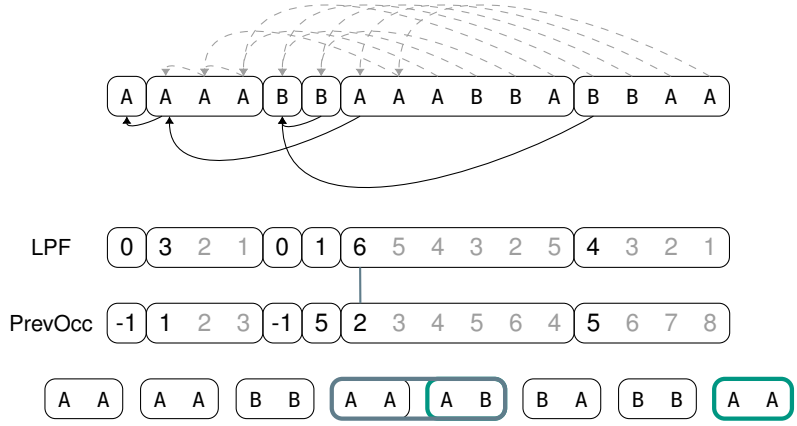

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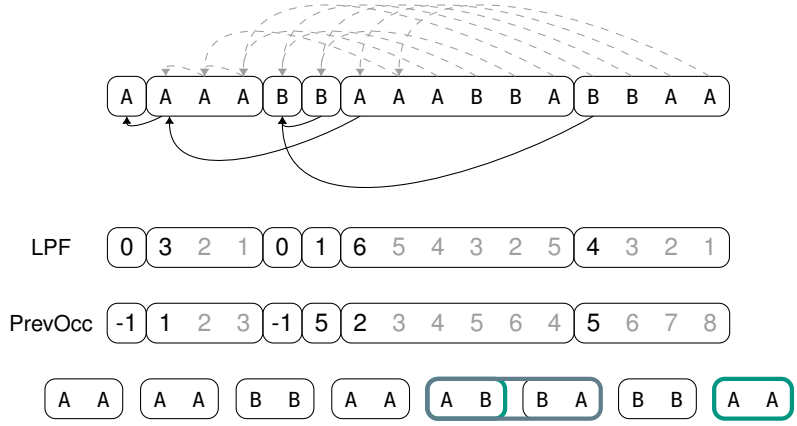

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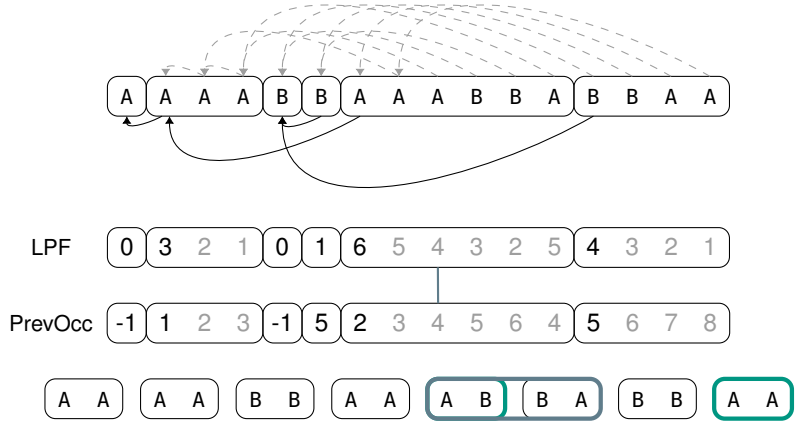

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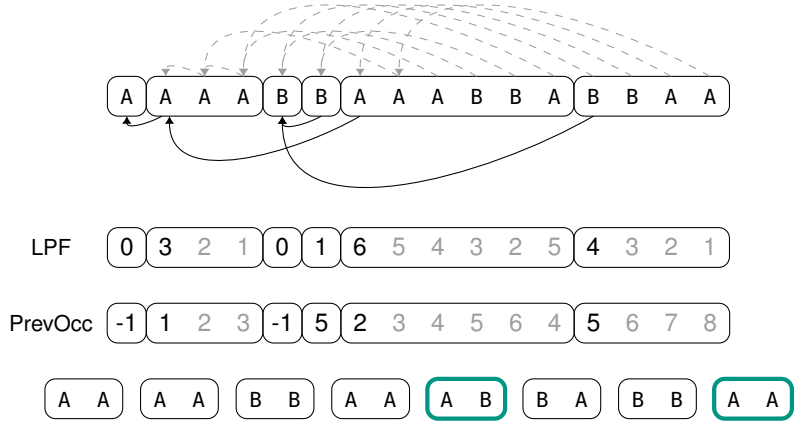

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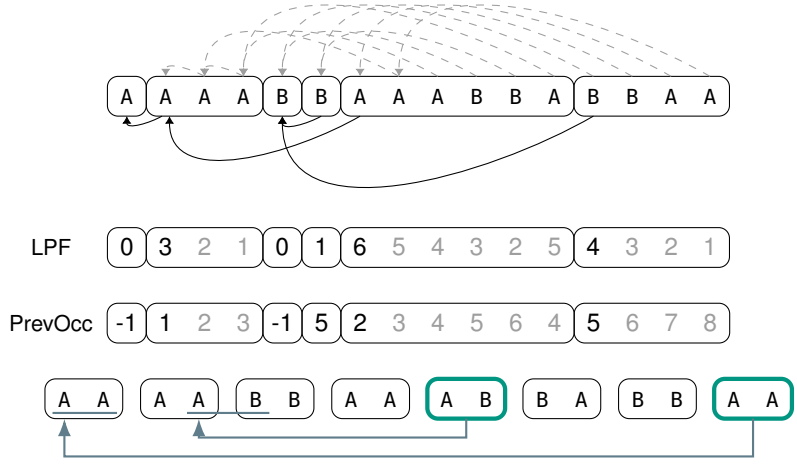

# Experimental Evaluation

- highly tuned implementation
- tree consists only of bit and compact vectors
- tuning parameter
  - degree root $s = \{1, z\}$ (only we have $s = z$)
  - degree other nodes $\tau = \{2, 4, 8, 16\}$
  - number characters in leaves $b = \{2, 4, 8, 16\}$

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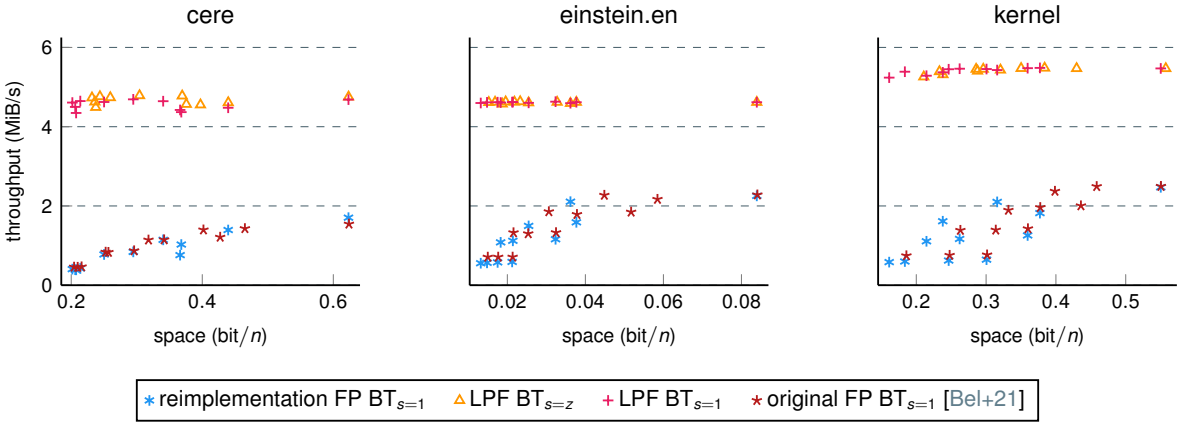
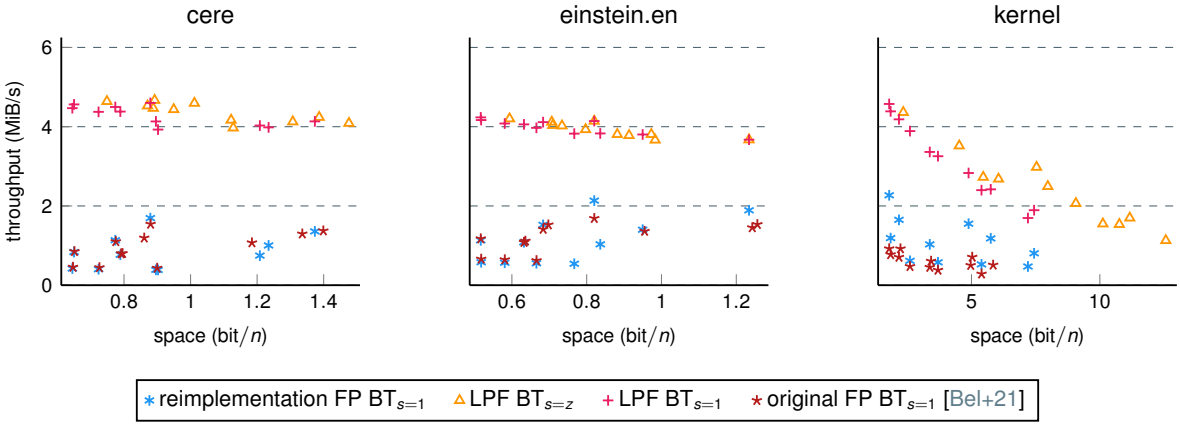
- original FP BT [Bel+21]
- our reimplementation of the original FP BT
- our LPF BT construction with $s = 1$ and $s = z$
- dynamic programming variants
- parallelization
- no comparison with wavelet trees (faster)

- repetitive instances from P&C corpus
- non-repetitive instances from P&C corpus

# Highly Repetitive Inputs (Access Only)



# Highly Repetitive Inputs (with Rank and Select Support)



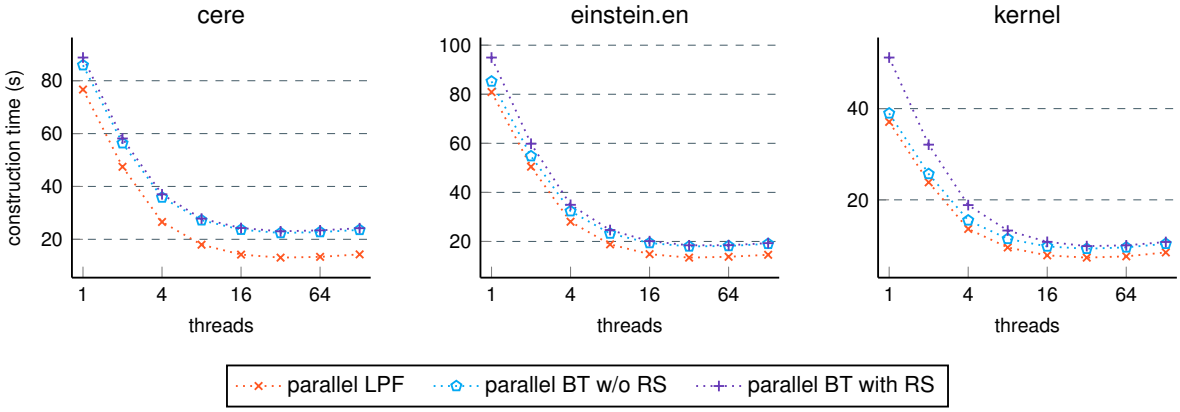
# Conclusion and Future Work

- fastest block tree construction algorithm
- first parallel block tree construction
- works in practice for non-repetitive inputs

- better scaling parallel construction
- better marking rules (less pruning)


This project has received funding from the European Research Council (ERC) under the European Union's Horizon 2020 research and innovation programme (grant agreement No. 882500).

# Parallel Construction (Strong Scaling)



# Bibliography

[Bel+21] Djamal Belazzougui, Manuel Cáceres, Travis Gagie, Pawel Gawrychowski, Juha Kärkkäinen, Gonzalo Navarro, Alberto Ordóñez Pereira, Simon J. Puglisi, and Yasuo Tabei. “Block Trees”. In: *J. Comput. Syst. Sci.* 117 (2021), pages 1–22. DOI: 10.1016/j.jcss.2020.11.002.